

Lesson 3.2 Exercises: Consensus Mechanisms and Blockchain Architecture

Module 3: The Trust Problem

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Digital Finance — BSc Course

Exercise 1: Byzantine Generals and Real Networks

Scenario: A consortium of 7 central banks is building a shared settlement system. Intelligence reports suggest that up to 2 of the banks' IT systems may be compromised by state-sponsored hackers who could inject fraudulent settlement messages.

Tasks:

- a Using the BFT bound $n > 3f$, determine whether the system can tolerate 2 Byzantine nodes out of 7. Show your calculation.
- b If the consortium grows to 10 banks, what is the maximum number of compromised nodes the system can tolerate?
- c Explain in plain language why the $\frac{1}{3}$ bound exists. Why can't a BFT system tolerate, say, 49% faulty nodes the way Proof of Work can tolerate up to 49% of hash power?
- d A junior engineer proposes: "If we suspect Bank X is compromised, let's just remove them from the validator set." Evaluate this proposal — what are the benefits and risks?

Difficulty: Introductory — tests understanding of BFT bounds and their practical meaning.

Exercise 2: Bitcoin Mining Profitability Analysis

Scenario: You are evaluating a Bitcoin mining investment. Use these parameters:

Parameter	Value
Mining rig hash rate	100 TH/s (terahashes per second)
Rig power consumption	3,000 W
Electricity cost	\$0.06 per kWh
Current block reward	3.125 BTC
Network hash rate	600 EH/s (6×10^{20} H/s)
BTC price	\$60,000
Blocks per day	144

Tasks:

- Calculate your rig's share of the total network hash rate.
- Calculate your expected daily BTC earnings (share \times total daily rewards).
- Calculate your daily electricity cost (power \times hours \times rate).
- Is the operation profitable? What BTC price would make it break-even?
- After the next halving (reward drops to 1.5625 BTC), recalculate. Is the operation still viable?

Difficulty: Intermediate — requires arithmetic and economic reasoning.

Exercise 3: Simulating Difficulty Adjustment

Scenario: A new PoW blockchain targets 5-minute block times and adjusts difficulty every 100 blocks. Use the data below:

Period	Actual Time for 100 Blocks	Difficulty
1	500 min (on target)	1,000
2	400 min (too fast)	?
3	?	?

Tasks:

- Calculate the new difficulty for Period 2 using: $\text{new} = \text{old} \times \frac{\text{target_time}}{\text{actual_time}}$.
- With the Period 2 difficulty, the network produces 100 blocks in 600 minutes. Calculate the Period 3 difficulty.
- Plot the three difficulty values. Is the system converging toward the 5-minute target?
- What happens if 50% of miners suddenly leave the network (e.g., due to a government ban) before the next adjustment? How long would blocks take on average? How does this compare to Bitcoin's experience after China's 2021 ban?

Difficulty: Intermediate — requires iterative calculation and critical analysis.

Exercise 4: Analyzing a Blockchain Fork

Scenario: At block height 500, two miners (A and B) simultaneously find valid blocks. The network splits:

- 55% of hash power builds on Miner A's block
- 45% of hash power builds on Miner B's block

Assume average block time is 10 minutes and the fork is resolved by the longest chain rule.

Tasks:

- What is the probability that Miner A's chain is extended first (i.e., one of the 55% miners finds block 501 before a 45% miner does)?
- After 3 blocks are built on Miner A's chain and only 2 on Miner B's chain, what happens to the transactions that were in Miner B's blocks but not in Miner A's?
- A merchant accepted a 10 BTC payment in Miner B's block 500. After the fork resolves in favor of Miner A's chain, has the merchant been paid? What should the merchant have done differently?
- Calculate the probability that Miner B's chain catches up after falling behind by 1 block, 2 blocks, and 3 blocks. Use the formula $P = (q/p)^k$ where $p = 0.55$, $q = 0.45$, $k = \text{deficit}$.

Difficulty: Advanced — requires probability reasoning and understanding of fork mechanics.

Exercise 5: Comparing Attack Costs

Scenario: You are a security consultant comparing the cost of attacking two networks:

Network X (PoW):

- Total hash rate: 500 EH/s
- Cost of hash power: \$0.05 per TH/s per day (rental)
- A 51% attack requires controlling ≥ 250 EH/s

Network Y (PoS):

- Total staked: 30 million ETH at \$2,000/ETH
- A 34% attack requires controlling ≥ 10.2 million ETH

Tasks:

- Calculate the daily cost of a 51% attack on Network X.
- Calculate the capital required to acquire 34% of Network Y's stake on the open market. Why would the actual cost be even higher than this calculation suggests?
- Compare the two attack costs. Which network is more expensive to attack?
- In Network Y, after a successful attack, the attacker's stake is slashed. What is the attacker's total loss? Does PoW have an equivalent "slashing" mechanism?

Difficulty: Advanced — requires quantitative security analysis across paradigms.

Exercise 6: Designing for the Trilemma

Scenario: Three blockchain projects approach you for advice. Each has a specific priority and asks how to configure their system:

- 1 **Project A — Global remittance platform:** Must handle 10,000 TPS with sub-5-second finality. Willing to sacrifice decentralization.
- 2 **Project B — Censorship-resistant digital gold:** Maximum decentralization and security. Throughput is secondary.
- 3 **Project C — Government digital identity system:** Must be highly secure and scalable, but can be centrally managed.

Tasks:

- a For each project, identify which two trilemma properties are prioritized and which is sacrificed.
- b Recommend a specific consensus mechanism for each project. Justify your choice.
- c For Project A, the team later demands “but we also want full decentralization.” Explain why this is infeasible with current technology and propose a Layer 2 compromise.
- d Draw a trilemma triangle and plot all three projects on it, showing their position.

Difficulty: Advanced — requires synthesis of trilemma trade-offs and mechanism selection.

Exercise 7: Evaluating a DeFi Protocol's Consensus Choice

Scenario: A DeFi lending protocol operates on three different blockchains and must choose which to prioritize:

	Chain P (PoW)	Chain Q (PoS)	Chain R (BFT)
Finality	~60 min	~13 min	~3 sec
TPS	7	30	5,000
Validators	10,000+	800,000+	21
Gas fees	\$15–\$50	\$0.50–\$5	\$0.001

Tasks:

- a. For a **liquidation engine** that must execute within seconds to prevent bad debt, which chain is best? Justify.
- b. For a **long-term savings vault** where security matters most and speed is irrelevant, which chain is best?
- c. Chain R has only 21 validators. A competitor bribes 8 of them. What attack is now possible? Would this attack be feasible on Chain Q?
- d. The protocol decides to deploy on all three chains. Propose a strategy that uses each chain's strengths (e.g., liquidations on R, long-term storage on P). Justify.

Difficulty: Advanced — requires multi-criteria evaluation and practical system design.

Exercise 8: Design a Cross-Border Payment Network

Scenario: The Central Bank of Euroland commissions you to design a blockchain-based cross-border payment system connecting 50 banks across 20 countries. Requirements:

- Finality within 10 seconds (irreversible settlements)
- Throughput: 500 transactions per second at peak
- Must tolerate up to 5 compromised banks
- Full auditability by regulators
- No single country should control the network

Tasks:

- a) Select a consensus mechanism. Justify why PoW, PoS, or BFT is most appropriate. Verify that the BFT bound $n > 3f$ is satisfied.
- b) Design the validator set: how many validators, how are they selected, what are the hardware requirements?
- c) Define the finality rule: when exactly is a transaction considered settled?
- d) Describe what happens when 6 banks (exceeding the tolerance threshold) go offline during a network partition. How does the system recover?
- e) Identify two risks that your consensus choice introduces and propose mitigations.

Difficulty: Advanced–Integrative — combines all lesson concepts in a complete system design.

Exercise 1:

- (a) $n > 3f$: $7 > 3 \times 2 = 6$. Yes, $7 > 6$, so the system can tolerate 2 Byzantine nodes.
- (b) $10 > 3f \Rightarrow f < 3.33$. Maximum $f = 3$ compromised nodes.
- (c) With $\frac{1}{3}$ faulty nodes, a Byzantine node can send "attack" to half the honest nodes and "retreat" to the other half, perfectly splitting the honest majority. In PoW, the chain is selected by cumulative work, not votes — an attacker with 49% hash power still produces fewer blocks on average, so the honest chain wins probabilistically.
- (d) Benefits: reduces attack surface. Risks: false accusations could remove honest banks; the "remove" mechanism itself could be weaponized by a Byzantine coalition to eject honest validators and gain majority.

Exercise 2:

- (a) Share = $100 \text{ TH/s} \div 600,000,000 \text{ TH/s} = 1.67 \times 10^{-7}$ (0.0000167%).
- (b) Daily BTC = $144 \times 3.125 \times 1.67 \times 10^{-7} = 0.0000751 \text{ BTC} = \4.50 .
- (c) Daily electricity = $3.0 \text{ kW} \times 24 \text{ h} \times \$0.06 = \$4.32$.
- (d) Profit = $\$4.50 - \$4.32 = \$0.18/\text{day}$. Barely profitable. Break-even at BTC $\approx \$57,600$.
- (e) After halving: daily BTC revenue halves to $\$2.25$. Cost remains $\$4.32$. Loss of $\$2.07/\text{day}$ — no longer viable without cheaper power or better hardware.

Answer Key (continued)

Exercise 3:

- (a) Target: 500 min. Actual: 400 min. New difficulty = $1,000 \times \frac{500}{400} = 1,250$.
- (b) Actual: 600 min. New difficulty = $1,250 \times \frac{500}{600} \approx 1,042$.
- (c) Difficulty: $1,000 \rightarrow 1,250 \rightarrow 1,042$. The system oscillates but converges toward equilibrium.
- (d) If 50% of miners leave, block time roughly doubles (to ~ 10 min) until the next adjustment. After China's 2021 ban, Bitcoin's hash rate dropped $\sim 50\%$ and block times temporarily increased to ~ 14 – 18 minutes before difficulty adjusted downward.

Exercise 4:

- (a) $P(\text{A extends first}) = 0.55$ (55%). Miner selection is proportional to hash rate.
- (b) Transactions unique to B's orphaned blocks return to the mempool and can be included in future blocks on A's chain. They are not lost, just delayed.
- (c) The merchant has NOT been paid — the transaction was in an orphaned block. The merchant should have waited for 6 confirmations on the canonical chain before releasing goods.
- (d) $P(1\text{-block}) = 0.45/0.55 = 0.818$. $P(2\text{-block}) = (0.818)^2 = 0.669$. $P(3\text{-block}) = (0.818)^3 = 0.548$. Each additional block makes catch-up exponentially less likely.

Answer Key (continued)

Exercise 5:

- (a) $250 \text{ EH/s} = 250,000,000 \text{ TH/s} \times \$0.05/\text{day} = \$12,500,000/\text{day}$.
- (b) $10.2\text{M ETH} \times \$2,000 = \$20.4 \text{ billion}$. Actual cost higher due to market slippage: buying 10.2M ETH would dramatically drive up the price.
- (c) PoS attack (**\$20.4B** capital) is orders of magnitude more expensive than PoW (**\$12.5M/day**). However, PoW cost is recurring; PoS capital is theoretically recoverable (if not slashed).
- (d) In PoS, the 10.2M ETH is slashed (destroyed). Total loss: $\sim \$20.4\text{B}$. PoW has no equivalent: the attacker's hardware still works, and electricity is already spent. PoS punishes the attacker's capital directly.

Exercise 6:

- (a) A: scalability + security, sacrifices decentralization. B: decentralization + security, sacrifices scalability. C: security + scalability, sacrifices decentralization.
- (b) A: Tendermint BFT or Avalanche (fast finality, high TPS, permissioned-friendly). B: PoW like Bitcoin (maximally decentralized, battle-tested). C: PBFT or permissioned PoS (known validators, high throughput).
- (c) No known system achieves all three at scale. Compromise: use a highly decentralized L1 (e.g., Ethereum) for security and settlement, with a high-throughput L2 rollup for execution.
- (d) Students draw triangle with A near the scalability-security edge, B near decentralization-security, C near security-scalability.

Answer Key (continued)

Exercise 7:

- (a) Chain R (BFT): 3-second finality and 5,000 TPS enable sub-second liquidation execution. Chain P's 60-minute finality would allow cascading bad debt. Chain Q's 13 minutes may be too slow for volatile markets.
- (b) Chain P (PoW): 10,000+ miners, decades of proven security, highest decentralization. For savings that sit for months/years, 60-minute finality is irrelevant; security and immutability are paramount.
- (c) 8 out of 21 validators exceeds the $\frac{1}{3}$ threshold ($7 < 21/3 = 7$, but $8 > 7$). The bribed coalition can halt finality or produce conflicting blocks. On Chain Q with 800,000 validators, bribing $\frac{1}{3}$ (267,000) is practically impossible.
- (d) Multi-chain strategy: liquidations and high-frequency operations on Chain R (speed); settlement and long-term treasury on Chain P (security); general DeFi operations on Chain Q (balance of cost, speed, decentralization). Cross-chain bridges pass settlement proofs from R and Q to P periodically.

Exercise 8:

- (a) BFT (e.g., Tendermint): 10-second finality + 500 TPS fits BFT's strengths. PoW too slow; PoS possible but BFT's deterministic finality is preferred for settlement. $n = 50 > 3 \times 5 = 15$. BFT bound satisfied.
- (b) 50 validators (one per bank). Geographically distributed across 20 countries. High-availability requirements: redundant hardware, 99.9% uptime SLA.
- (c) A transaction is final once $\frac{2}{3}$ of validators (34+) sign the commit message. Typically 2–3 seconds.
- (d) 6 offline $>$ 5 tolerance. Finality halts (only 44 of 50 online; need 34, so finality still works if $44 \geq 34$). Actually 44 validators can still finalize. But if 17+ go offline, finality halts. Recovery: inactivity penalties or manual re-admission once banks come back online.
- (e) Risk 1: Validator collusion (geopolitical blocs). Mitigation: require validators from different jurisdictions in every quorum. Risk 2: Single point of failure in leader selection. Mitigation: rotate leaders every block using verifiable random functions.